

Advanced Algorithms VI (Fall 2020)

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In this lecture, we prove the Johnson-Lindenstrauss Lemma, which is an important algorithmic application of concentration inequalities. To this end, we introduce the notion of sub-Gaussian random variables and the Bernstein's inequality.

1 Johnson-Lindenstrauss Lemma

Given a collection of points in a high-dimensional space, one might try to map the points into a low-dimensional space without distorting the relative distances between points very much. This problem is called *metric embedding* in literature. The Johnson-Lindenstrauss lemma states that such an embedding exists in certain case.

Lemma 1 (Johnson-Lindenstrauss Lemma). *Let S be a collection of points set that $S \subseteq \mathbb{R}^D$. Then for every $\varepsilon \in (0, 1)$, there exists a projection $f : \mathbb{R}^D \rightarrow \mathbb{R}^d$ with $d = O(\log(|S|/\varepsilon^2))$ such that $\forall x, y \in S, x \neq y$, it holds that*

$$1 - \varepsilon \leq \frac{\|f(x) - f(y)\|}{\|x - y\|} \leq 1 + \varepsilon \quad (1)$$

Note that in the statement of JL lemma, the dimension d is irrelevant to D . So D can be arbitrary large or even infinite. It is surprising that a *random* linear projection from \mathbb{R}^D to \mathbb{R}^d satisfies the requirement of JL with high probability. To see this, suppose we define a matrix $A = (a_{ij})_{\substack{1 \leq i \leq D \\ 1 \leq j \leq d}} \in \mathbb{R}^{D \times d}$ where each a_{ij} is chosen from $\{-1, 1\}$ uniformly at random. Then for any vector $u \in \mathbb{R}^D$, $Au \in \mathbb{R}^d$ satisfies

$$\mathbf{E} [\|Au\|^2] = \mathbf{E} \left[\sum_{i=1}^d (Au)_i^2 \right] = \sum_{i=1}^d \mathbf{E} \left[\left(\sum_{j=1}^D a_{ij} \cdot u_j \right)^2 \right] = d \cdot \|u\|^2.$$

Therefore, if we choose $f = \frac{A}{\sqrt{d}}$, then $\mathbf{E} [\|f(u)\|^2] = \|u\|^2$. This implies that for any $x, y \in S$, $\mathbf{E} [\|f(x) - f(y)\|^2] = \mathbf{E} [\|f(x - y)\|^2] = \mathbf{E} [\|x - y\|^2]$. Hence to establish (1), we only need to prove that $\|f(u)\|^2$ is well-concentrated to its expectation, namely

$$\Pr \left[1 - \varepsilon \leq \frac{\|f(u)\|^2}{\|u\|^2} \leq 1 + \varepsilon \right] \geq 1 - \delta$$

for appropriate δ . Since f is linear, we can assume without loss of generality that $\|u\| = 1$.

For every $i = 1, \dots, d$, we let $Z_i = \sum_{j=1}^D a_{ij}u_j$. Then $\|f(u)\|^2 = \frac{1}{d} \cdot \sum_{i=1}^d Z_i^2$. Now we can express our objective as the sum of d independent random variables, so to obtain a concentration result, we might

try to apply Chernoff-typed inequalities. However, it seems that we cannot directly apply the Hoeffding inequality here, since $Z_i^2 = \left(\sum_{j=1}^D a_{ij}u_j\right)^2$ can be unbounded. Therefore, we need some new tools to tackle random variables of this form.

2 Sub-Gaussian Random Variables

Recall the proof of the Chernoff bounds and the Hoeffding inequality. The key to establish these inequalities is an upper bound on the moment generating functions $\mathbf{E}\left[e^{\alpha X}\right]$. We abstract the property and introduce the notion of *sub-Gaussian* random variables.

Definition 2. A random variables X with $\mathbf{E}[X] = 0$ is called *sub-Gaussian with variance factor v* , denoted as $X \in \mathcal{G}(v)$, if

$$\mathbf{E}\left[e^{\alpha X}\right] \leq e^{\frac{\alpha^2}{2}v} \quad \text{for every } \alpha \in \mathbb{R}.$$

The name *sub-Gaussian* comes from the fact that for a Gaussian random variables $X \sim N(0, v)$, it holds that $\mathbf{E}\left[e^{\alpha X}\right] = e^{\frac{\alpha^2}{2}v}$.

The moment generating function of a random variables X is closely related to its k -the moment for all $k \in \mathbb{N}$. The following theorem clarifies the relationship, and interested readers can refer to [1, Chapter 2] for more on this.

Theorem 3. Let X be a random variable with $\mathbf{E}[X] = 0$

(1) If $X \in \mathcal{G}(v)$, then for every integer $j \geq 1$, $\mathbf{E}\left[X^{2j}\right] \leq 2^{j+1}j!v^j$

(2) If for some positive constant v and for every integer $j \geq 1$, $\mathbf{E}\left[X^{2j}\right] \leq j!v^j$, then $X \in \mathcal{G}(4v)$

Proof. We first prove (1). We can assume without loss of generality that X is a continuous random variable. Assume $X \in \mathcal{G}(v)$, then we have

$$\begin{aligned} \mathbf{E}\left[X^{2j}\right] &= \int_0^\infty \Pr\left[|X|^{2j} \geq x\right] dx \\ &= \int_0^\infty \Pr\left[|X| \geq x^{\frac{1}{2j}}\right] dx \\ &= 2j \int_0^\infty \Pr\left[|X| \geq z\right] z^{2j-1} dz \quad (z = x^{\frac{1}{2j}}). \end{aligned}$$

For any $\alpha > 0$, we have

$$\Pr[X > z] = \Pr\left[e^{\alpha X} > e^{\alpha z}\right] \leq \frac{\mathbf{E}\left[e^{\alpha X}\right]}{e^{\alpha z}} \leq e^{\frac{\alpha^2}{2}v - \alpha z}.$$

We choose $\alpha = \frac{z}{v}$ and obtain $\Pr[X > z] \leq e^{-\frac{z^2}{2v}}$. Similarly, we can obtain $\Pr[X < -z] \leq e^{-\frac{z^2}{2v}}$. Then

$$\begin{aligned} \mathbf{E}\left[X^{2j}\right] &\leq 4j \int_0^\infty z^{2j-1} e^{-\frac{z^2}{2v}} dz \\ &= 4j \int_0^\infty (2vt)^{j-\frac{1}{2}} e^{-t} d(2vt)^{\frac{1}{2}} \quad (t = \frac{z^2}{2v}) \\ &= 2j(2v)^j \int_0^\infty t^{j-1} e^{-t} dt \\ &= 2^{j+1}j!v^j. \end{aligned}$$

We proceed to prove (2). Assume $\mathbf{E} [X^{2j}] \leq j!v^j$. To get rid of the odd moments of X , we introduce an independent random variable X' who follows the same distribution as X . Then by symmetry of $X - X'$ we have

$$\mathbf{E} [e^{\alpha X}] \mathbf{E} [e^{-\alpha X'}] = \mathbf{E} [e^{\alpha X} e^{-\alpha X'}] = \mathbf{E} [e^{\alpha(X-X')}] = \sum_{j=0}^{\infty} \frac{\alpha^j \cdot \mathbf{E} [(X - X')^j]}{j!}.$$

For odd j , we have

$$\begin{aligned} \mathbf{E} [(X - X')^j] &= \sum_{k=0}^j \binom{j}{k} \mathbf{E} [X^k] \mathbf{E} [(-X')^{j-k}] \\ &= \sum_{k=0}^j (-1)^{j-k} \cdot \binom{j}{k} \mathbf{E} [X^k] \mathbf{E} [X^{j-k}] \\ &= \sum_{k=0}^{\lfloor j/2 \rfloor} \binom{j}{k} \left((-1)^k + (-1)^{j-k} \right) \mathbf{E} [X^k] \mathbf{E} [X^{j-k}] \\ &= 0. \end{aligned}$$

Therefore,

$$\mathbf{E} [e^{\alpha X}] \mathbf{E} [e^{-\alpha X'}] = \sum_{j=0}^{\infty} \frac{(\alpha)^{2j} \mathbf{E} [(X - X')^{2j}]}{(2j)!}. \quad (2)$$

Since the function $X \rightarrow X^{2j}$ is convex, Jensen's inequality yields

$$(X + (-X'))^{2j} = 2^{2j} \left(\frac{1}{2}X + \frac{1}{2}(-X') \right)^{2j} \leq 2^{2j} \left(\frac{1}{2}X^{2j} + \frac{1}{2}X'^{2j} \right) = 2^{2j-1}(X^{2j} + X'^{2j}).$$

So

$$\mathbf{E} [(X - X')^{2j}] \leq 2^{2j-1} \mathbf{E} [X^{2j} + X'^{2j}] = 2^{2j} \mathbf{E} [X^{2j}].$$

Since

$$\frac{(2j)!}{j!} = \prod_{k=1}^j (j+k) \geq \prod_{k=1}^j 2k = 2^j j!,$$

we have

$$(2) = \sum_{j=0}^{\infty} \frac{(\alpha)^{2j} \mathbf{E} [(X - X')^{2j}]}{(2j)!} \leq \sum_{j=0}^{\infty} \frac{(\alpha)^{2j} 2^{2j} \mathbf{E} [X^{2j}]}{(2j)!} \leq \sum_{j=0}^{\infty} \frac{(\alpha)^{2j} 2^{2j} e^j j!}{(2j)!} \leq \sum_{j=0}^{\infty} \frac{\alpha^{2j} v^j 2^j}{j!}.$$

Moreover, again by Jensen's inequality, we know that

$$\mathbf{E} [e^{-\alpha X'}] \geq e^{-\alpha \mathbf{E}[X']} = 1.$$

So

$$\mathbf{E} [e^{\alpha X}] \leq \sum_{j=0}^{\infty} \frac{\alpha^{2j} v^j 2^j}{j!} = e^{2\alpha^2 v}.$$

That is, $X \in \mathcal{G}(4v)$. □

Given bounds on all moments, we have the following more general concentration inequality, known as *Bernstein's inequality*.

Theorem 4 (Bernstein's inequality). *Let X_1, \dots, X_n be independent real-valued random variables. Assume that there exist positive numbers a and b such that*

$$(1) \sum_{i=1}^n \mathbf{E} [X_i^2] \leq a$$

$$(2) \sum_{i=1}^n \mathbf{E} [X_i^j] \leq \frac{j!}{2} ab^{j-2} \text{ for all integers } q \geq 3$$

Define $X = \sum_{i=1}^n X_i$, $S = X - \mathbf{E}[X]$, then for all $t > 0$, we have

$$\Pr [S \geq \sqrt{2at} + bt] \leq e^{-t}.$$

See [1] for a proof of the theorem.

3 Proof of Johnson-Lindenstrauss Lemma

We are now ready to prove Lemma 1. Following the discussion in Section 1, for every $i = 1, \dots, d$, we have

$$\mathbf{E} [e^{\alpha Z_i}] = \mathbf{E} [e^{\alpha \sum_{j=1}^D a_{ij} u_j}] = \prod_{j=1}^D \mathbf{E} [e^{\alpha a_{ij} u_j}] = \prod_{j=1}^D \left(\frac{1}{2} (e^{\alpha u_j} + e^{-\alpha u_j}) \right).$$

Since

$$\frac{1}{2}(e^\lambda + e^{-\lambda}) = \sum_{j=0}^{\infty} \frac{\lambda^{2j}}{(2j)!} \leq \sum_{j=0}^{\infty} \frac{(\lambda)^{2j}}{2^j j!} = e^{\frac{\lambda^2}{2}},$$

we have

$$\mathbf{E} [e^{\alpha Z_i}] \leq \prod_{j=1}^D e^{\frac{\alpha^2 u_j^2}{2}} = e^{\frac{\alpha^2}{2}}.$$

Therefore, $Z_i \in \mathcal{G}(1)$. Let $Y_i = Z_i^2$. According to Theorem 3, we can obtain bounds on moments of Z_i and Y_i :

$$\forall j \geq 1 : \mathbf{E} [Z_i^{2j}] \leq 2^{j+1} \cdot j!;$$

$$\forall j \geq 1 : \mathbf{E} [Y_i^j] \leq 2^{j+1} \cdot j! \leq 4^j \cdot j!;$$

$$\forall j \geq 1 : \mathbf{E} [Y_i^{2j}] \leq 2^{2j+1} \cdot (2j)! \leq 2^{3j+1} \cdot j!.$$

Finally we have

$$\Pr [\|f(u)\|^2 - 1 > \varepsilon] = \Pr \left[\sum_{i=1}^d \frac{1}{d} Y_i - 1 > \varepsilon \right] = \Pr \left[\sum_{i=1}^d (Y_i - 1) > \varepsilon d \right].$$

We can let $a = 16d$ and $b = 4$ in the Bernstein's inequality (Theorem 4), which gives

$$\Pr \left[\sum_{i=1}^d Y_i - 1 \geq 4\sqrt{2dt} + 4t \right] \leq e^{-t}.$$

Applying union bound for every pair of $x, y \in S$, it suffices to let $n^2 \cdot e^{-t} \leq \delta$. So we let $4\sqrt{2dt} + 4t = d\varepsilon$ where $t = \log \frac{n^2}{\delta}$. This requires $d = \Theta \left(\frac{1}{\varepsilon^2} \log \frac{n}{\sqrt{\delta}} \right)$.

References

- [1] S. BOUCHERON, G. LUGOSI, AND P. MASSART, *Concentration inequalities: A nonasymptotic theory of independence*, 2013. 2, 4