

[CS3958: Lecture 14] Graph Expansion(Cont'), Cheeger's Inequality

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1 Graph Expansion (Cont')

Expansion can be defined on any weighted undirected graph, not only for Markov chains: Let $G = (V, E)$ be a weighted undirected graph with a weight function $w(i, j) > 0$ for each edge $\{i, j\} \in E$. Then we define the expansion as

$$\Phi(S, \bar{S}) = \frac{\sum_{i \in S, j \in \bar{S}} w(i, j)}{\sum_{i, j \in S} w(i, j)}.$$

Note this definition is consistent with that of Markov chains. If we let $P(i, j) = \frac{w(i, j)}{\sum_j w(i, j)}$ i.e. P is a natural random walk on G , then we have $\pi(i) \sim \sum_j w(i, j)$ and $\pi(i)P(i, j) = \pi(j)P(j, i)$. Therefore, the Markov chain P can imply some results on the expansion of G .

1.1 Applications for Sampling Colorings

Assume we want to sample from all proper $[q]$ -colorings on $G = ([n], E)$ with maximum degree Δ . The Markov chain is

- Pick $v \in [n]$ and $c \in [q]$ uniformly at random.
- Recolor v with c if possible.

Recall that we proved $\tau_{mix}(\epsilon) \leq qn \log \frac{n}{\epsilon}$ when $q > 4\Delta$. Now we want to argue that when q is rather small, the expansion is large for some special graph. Consider the case when G is a star and 1 is the vertex at the center. Let Z be the number of all proper colorings on G , and S be the set of proper colorings that the color of 1 is 1. Then we have

$$\begin{aligned} Q(S, \bar{S}) &= \sum_{i \in S, j \in \bar{S}} \pi(i)P(i, j) \\ &= (q-1)(q-2)^{n-1} \frac{1}{Z \cdot nq}. \end{aligned}$$

Since $|S| = (q-1)^{n-1}$, we have

$$\Phi(S) = \frac{Q(S, \bar{S})}{\pi(S)} = \frac{(q-2)^{n-1} \frac{1}{nq}}{(q-1)^{n-2} \frac{1}{nq}} = \frac{q-1}{nq} \left(1 - \frac{1}{q-1}\right)^{n-2} \leq \frac{1}{n} \exp\left(-\frac{n-2}{q-1}\right).$$

Therefore, $\tau_{mix} = \Omega\left(n \cdot \exp\left(\frac{q-1}{n-2}\right)\right)$, which means that when $q = o\left(\frac{n}{\log n}\right)$, τ_{mix} is $n^{\omega(1)}$.

Review: Different views of analyzing mixing time/rate of convergence of Markov Chains.

- Probabilistic view ~ Coupling;
- Algebraic view ~ Spectrum; (Algebraic Graph Theory)
- Geometric view ~ Expansion. (KLS Conjecture)

If we want to upperbound $\Phi(P)$, we need to argue that for any S such that $\pi(S) \leq \frac{1}{2}$, $\Phi(S, \bar{S})$ has an upper bound.

The problem to find a cut S, \bar{S} with maximum expansion is dual with multi-commodity flow problem. If you are interested in this topic, search for "canonical paths" or "multi-commodity flow".

2 Cheeger's Inequality

Sometimes it is more convenient to work with $L = I - P$, the *Laplacian* of P .

Then

$$L = \sum_{i=1}^n (1 - \lambda_i) \mathbf{v}_i \mathbf{v}_i^\top \Pi.$$

For every $i = 1, 2, \dots, n$, we use γ_i denote $1 - \lambda_i$. Then $0 = \gamma_1 \leq \gamma_2 \leq \dots \leq \gamma_n \leq 2$ are the eigenvalues of L .

The Cheeger's inequality is

$$\frac{\gamma_2}{2} \leq \Phi(P) \leq \sqrt{2\gamma_2}.$$

There are high-order Cheeger's inequalities indicating the relation between graph expansion and $\lambda_3, \lambda_4, \dots$

2.1 Proof of $\gamma_2 \leq 2\Phi(P)$

Recall that

$$\gamma_2 = \min_{2-\dim V \subseteq \mathbb{R}^n} \max_{\mathbf{x} \in V \setminus \{0\}} R_L(\mathbf{x}).$$

Therefore, in order to prove an upper bound for γ_2 , it suffices to construct some 2-dimensional space V such that any nonzero $\mathbf{x} \in V$ has small $R_L(\mathbf{x})$.

Suppose $\Phi(P) = \Phi(S)$ for some $S \subseteq V$. Let $\mathbf{1}_S$ and $\mathbf{1}_{\bar{S}}$ be the indicator vector of S and its complement \bar{S} respectively. Consider the space $V = \text{span}(\mathbf{1}_S, \mathbf{1}_{\bar{S}})$. Then every $\mathbf{x} \in V$ can be written as $\mathbf{x} = a\mathbf{1}_S + b\mathbf{1}_{\bar{S}}$ for some $a, b \in \mathbb{R}$. We have

$$\begin{aligned} R_L(a\mathbf{1}_S) &= \frac{\sum_{i \in S, j \in \bar{S}} \pi(i)P(i, j)}{\pi(S)} \\ &= \frac{\sum_{i \in S, j \in \bar{S}} \Pr[X_t = i \wedge X_{t+1} = j]}{\pi(S)} \\ &= \frac{\Pr[X_t \in S \wedge X_{t+1} \in \bar{S}]}{\pi(S)} \\ &= \Phi(S). \end{aligned}$$

Similarly, we have $R_L(b\mathbf{1}_{\bar{S}}) = \Phi(\bar{S})$.

The inequality then follows from the following proposition:

Proposition 1 *If at least one of \mathbf{x} and \mathbf{y} is not zero, then $R_L(\mathbf{x} + \mathbf{y}) \leq 2 \max\{R_L(\mathbf{x}), R_L(\mathbf{y})\}$.*

Assume $\mathbf{x} = \sum_{i=1}^n a_i \mathbf{v}_i$ and $\mathbf{y} = \sum_{i=1}^n b_i \mathbf{v}_i$. Then

$$\begin{aligned} R_L(\mathbf{x} + \mathbf{y}) &= \frac{\langle \mathbf{x} + \mathbf{y}, L(\mathbf{x} + \mathbf{y}) \rangle_\Pi}{\langle \mathbf{x} + \mathbf{y}, \mathbf{x} + \mathbf{y} \rangle_\Pi} \\ &= \frac{\sum_{i=1}^n (a_i + b_i)^2 \lambda_i}{\sum_{i=1}^n (a_i + b_i)^2} \\ &\leq \frac{2 \sum_{i=1}^n (a_i^2 + b_i^2) \lambda_i}{\sum_{i=1}^n a_i^2 + b_i^2} \\ &\leq 2 \cdot \max\{R_L(\mathbf{x}), R_L(\mathbf{y})\}. \end{aligned}$$

2.2 Proof of $\Phi(P) \leq \sqrt{2\gamma_2}$

In order to prove an upper bound for $\Phi(P)$, we give an approximation algorithm to estimate $\Phi(P)$ and the upper bound is a consequence of the analysis of its performance. The algorithm is called *Fiedler's Algorithm*:
 Input Ω and $\mathbf{x} \in \mathbb{R}^\Omega$.

Our goal is to find (S, \bar{S}) such that $\Phi(S) \vee \Phi(\bar{S}) \leq \sqrt{2\gamma_2}$.

- Sort $\Omega = \{v_1, \dots, v_n\}$ according to \mathbf{x} (namely $x(v_1) \leq x(v_2) \leq \dots$);
- For every $i \in [n]$, let $S_i = \{v_1, v_2, \dots, v_i\}$;
- Return $\min_{i \in [n]} \Phi(S_i) \vee \Phi(\bar{S}_i)$.

We prove the following stronger theorem:

Theorem 2 For all $\mathbf{x} \perp \mathbf{1}$, let S be the set returned by Fiedler's algorithm on the input \mathbf{x} . Then

$$\Phi(S) \leq \sqrt{2R_L(\mathbf{x})}.$$

The Cheeger's inequality then follows by taking $\mathbf{x} = \mathbf{v}_2$.

For simplicity, we assume $\Omega = [n]$ and $\mathbf{x}(1) \leq \mathbf{x}(2) \leq \dots$ here. To prove the theorem, we first normalize the vector \mathbf{x} . Let

$$\ell \triangleq \min_{i \in [n]} \sum_{i=1}^{\ell} \pi(i) \geq \sum_{i=\ell+1}^n \pi(i),$$

and for every $i \in [n]$, let $y_i = x_i - x_\ell$. By the definition, $y_\ell = 0$ and $y_i \leq 0$ for all $i \leq \ell$, $y_i \geq 0$ for all $i \geq \ell$.

We have the following proposition:

Proposition 3 $R_L(\mathbf{x}) \geq R_L(\mathbf{y})$.

To see why it holds, note that

$$R_L(\mathbf{x}) = \frac{\sum_{i,j} \pi(i)P(i,j)(x_i^2 - x_i x_j)}{\langle \mathbf{x}, \mathbf{x} \rangle_\Pi} = \frac{\sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j)(x_i - x_j)^2}{\langle \mathbf{x}, \mathbf{x} \rangle_\Pi}.$$

Since $\mathbf{y} = \mathbf{x} - y_\ell \mathbf{1}$ is obtained from \mathbf{x} by subtracting a constant multiples of $\mathbf{1}$, this operation does not change the numerator and increase the denominator (because $\mathbf{x} \perp \mathbf{1}$). This can also be verified via direct calculation.

As a result, we only need to prove that $\Phi(S) \leq \sqrt{2R_L(\mathbf{y})}$. We prove by the *probabilistic method*. That is, we randomly choose some $t \in [y_1, y_n]$ (following a certain tailored density) and consider the expected expansions of $\Phi(S_t)$ and $\Phi(\bar{S}_t)$ where $S_t \triangleq \{i \in [n] \mid y_i \leq t\}$.

To this end, we can normalize \mathbf{y} by dividing some constant and assume without loss of generality that $y_1^2 + y_n^2 = 1$. We sample t with density $p(t) = 2|t|$.

Note that for every $t \in [y(1), y(n)]$,

$$\max \left\{ \Phi(S_t), \Phi(\bar{S}_t) \right\} = \frac{\sum_{i \in S_t, j \in \bar{S}_t} \pi(i)P(i,j)}{\min \left\{ \pi(S_t), \pi(\bar{S}_t) \right\}} =: \frac{A}{B}.$$

We calculate the expectations of the numerator and denominator respectively.

$$\begin{aligned}
\mathbf{E}[A] &= \sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) \Pr \left[i \in S_t, j \in \bar{S}_t \right] \\
&= \sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) \int_{y_j}^{y_i} 2|t| dt \\
&= \sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) \left(\operatorname{sgn}(y_j) \cdot y_j^2 - \operatorname{sgn}(y_i) \cdot y_i^2 \right) \\
&\leq \sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) (|y_i| + |y_j|) (y_j - y_i) \\
&= \sum_{\substack{i,j \in \Omega \\ i < j}} (\pi(i)P(i,j))^{\frac{1}{2}} (|y_i| + |y_j|) (\pi(i)P(i,j))^{\frac{1}{2}} (y_j - y_i) \\
&\stackrel{(\heartsuit)}{\leq} \left(\sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) (|y_i| + |y_j|)^2 \right)^{\frac{1}{2}} \cdot \left(\sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) (y_j - y_i)^2 \right)^{\frac{1}{2}} \\
&\leq \left(2 \sum_{\substack{i,j \in \Omega \\ i < j}} \pi(i)P(i,j) (|y_i|^2 + |y_j|^2) \right)^{\frac{1}{2}} \cdot \sqrt{\langle \mathbf{y}, L\mathbf{y} \rangle_{\Pi}} \\
&\leq \sqrt{2\langle \mathbf{y}, \mathbf{y} \rangle_{\Pi}} \cdot \sqrt{\langle \mathbf{y}, L\mathbf{y} \rangle_{\Pi}},
\end{aligned}$$

where (\heartsuit) is due to Cauchy-Schwarz.

On the other hand, we have

$$\begin{aligned}
\mathbf{E}[B] &= \mathbf{E} \left[\min \left\{ \pi(S_t), \pi(\bar{S}_t) \right\} \right] \\
&= \Pr[t < 0] \mathbf{E}[\pi(S_t) \mid t < 0] + \Pr[t \geq 0] \mathbf{E}[\pi(\bar{S}_t) \mid t \geq 0].
\end{aligned}$$

$$\langle \mathbf{y}, \mathbf{y} \rangle_{\Pi} = \sum_{i,j} \pi(i)P(i,j) y_i^2 \geq \sum_{i < j} \pi(i)P(i,j) (y_i^2 + y_j^2).$$

Note that

$$\begin{aligned}
\Pr [t < 0] \mathbf{E} [\pi(S_t) \mid t < 0] &= \Pr [t < 0] \cdot \mathbf{E} \left[\sum_{i=1}^n \pi(i) \cdot \mathbf{1}[i \in S_t] \mid t < 0 \right] \\
&= \Pr [t < 0] \cdot \sum_{i=1}^n \pi(i) \cdot \Pr [i \in S_t \mid t < 0] \\
&= \sum_{i=1}^n \pi(i) \cdot \Pr [i \in S_t \wedge t < 0] \\
&= \sum_{i=1}^n \pi(i) \Pr [y_i \leq t < 0] \\
&= \sum_{i \leq \ell} \pi(i) \cdot \int_{y_i}^0 2|t| dt \\
&= \sum_{i=1}^{\ell} \pi(i) y(i)^2.
\end{aligned}$$

Similarly

$$\Pr [t \geq 0] \mathbf{E} [\pi(\bar{S}_t) \mid t > 0] = \sum_{i=\ell+1}^n \pi(i) y(i)^2.$$

Therefore,

$$\mathbf{E} [B] = \sum_{i=1}^n \pi(i) y(i)^2 = \langle \mathbf{y}, \mathbf{y} \rangle_{\Pi}.$$

Now we know that

$$\frac{\mathbf{E} [A]}{\mathbf{E} [B]} \leq \frac{\sqrt{2\langle \mathbf{y}, \mathbf{y} \rangle_{\Pi}} \cdot \sqrt{\langle \mathbf{y}, L\mathbf{y} \rangle_{\Pi}}}{\langle \mathbf{y}, \mathbf{y} \rangle_{\Pi}} = \sqrt{2R_L(\mathbf{y})} \leq \sqrt{2R_L(\mathbf{x})}.$$

Moreover, for any r , we have

$$\frac{\mathbf{E} [A]}{\mathbf{E} [B]} \leq r \implies \mathbf{E} [A - rB] \leq 0 \implies \Pr \left[\frac{A}{B} \leq r \right] > 0.$$

The Cheeger's inequality is proved.